

[...]

(19a) **A farmer introduced a doctor to a priest. He told him that he already knew him.** Discourse(b) $\{(f,d,p) \mid f \text{ is a farmer, } d \text{ is a doctor, } p \text{ is a priest}\} =: R$ (c) $\{w,x,y,z\} \mid w \text{ told } x \text{ that } y \text{ already knew } z =: S$ (d) $R \oplus_{\{w=p,x=f,y=p,z=d\}} S = R \cap \{(f,d,p) \mid p \text{ told } f \text{ that } p \text{ already knew } d\}$ From Relations to Possibilities(20a) $R \subseteq U^n = \underbrace{U \times \dots \times U}_{n \text{ times}}$ set of n -tuples (Cartesian product)(20b) $B^A = \{f \mid f: A \rightarrow B\}$ set of functions from A to B (20c) $n = \{m \in \omega \mid m < n\} = \{0, \dots, n\}$ von Neumann ordinals(20d) $R \subseteq U^n = \{f: n \rightarrow U\}$ set of n -place sequences(20e) $U^n \cong \{f: X \rightarrow U\} \Leftrightarrow |X| = n$ (extensional) possibilitiesDiscourse RepresentationsBasic symbolspredicates (with intrinsic -arities): **P, Q, R ...**Variables (a.k.a. *discourse referents*): **x, y, z, ...**(binary) quantifiers: **∀, MOST, NO, ...**

auxiliary symbols: (,), [,], ,, :

Categories:Main category: *DRS*others: *Pred_n, Var, Cond, Quant*(Proper) Expressions (and their free variables)(i) $R \in \text{Pred}_n, x_1, \dots, x_n \in \text{Var} \Rightarrow R(x_1, \dots, x_n) \in \text{Cond};$

$$\text{Fr}(R(x_1, \dots, x_n)) = \{x_1, \dots, x_n\}$$

(ii) $\varphi, \psi \in \text{Cond} \Rightarrow [\varphi.\psi] \in \text{Cond};$

$$\text{Fr}([\varphi.\psi]) = \text{Fr}(\varphi) \cup \text{Fr}(\psi)$$

(iii) $X \subseteq \text{Var}, \varphi \in \text{Cond} \Rightarrow [X: \varphi] \in \text{DRS};$

$$\text{Fr}([X: \varphi]) = \text{Fr}(\varphi) \setminus X$$

(iv) $Q \in \text{Quant}, K, K' \in \text{DRS}, x \in \text{Fr}(K) \cap \text{Fr}(K') \Rightarrow [K \langle Qx \rangle K'] \in \text{Cond};$

$$\text{Fr}([K \langle Qx \rangle K']) = (\text{Fr}(K) \cup \text{Fr}(K')) \setminus \{x\}$$

Fairly standard interpretationModels: $\mathcal{M} = (\mathcal{U}_{\mathcal{M}}, \mathcal{F}_{\mathcal{M}})$ where: $\mathcal{U}_{\mathcal{M}} \neq \emptyset, \mathcal{F}_{\mathcal{M}}(\mathbf{R}) \subseteq \mathcal{U}_{\mathcal{M}}^n$ and $\mathcal{F}_{\mathcal{M}}(\mathbf{Q}) \subseteq \wp(\mathcal{U}_{\mathcal{M}}) \times \wp(\mathcal{U}_{\mathcal{M}})$ whenever $\mathbf{R} \in \text{Pred}_n$ and $\mathbf{Q} \in \text{Quant}$ Variable assignments g (based on model \mathcal{M}): $\text{Var} \rightsquigarrow \mathcal{U}_{\mathcal{M}}$ (partial functions)

Value $\llbracket \mathbf{A} \rrbracket^g$ of expression \mathbf{A} given \mathcal{M} and (suitable*) g :

$$(i) \quad \llbracket \mathbf{R}(x_1, \dots, x_n) \rrbracket^{\mathcal{M}, g} = \begin{cases} 1 & \text{iff } (g(x_1), \dots, g(x_n)) \in \mathcal{F}_{\mathcal{M}}(\mathbf{R}) \\ 0 & \text{iff } (g(x_1), \dots, g(x_n)) \notin \mathcal{F}_{\mathcal{M}}(\mathbf{R}) \end{cases}$$

(i*) where $\{x_1, \dots, x_n\} \subseteq \text{dom}(g)$

$$(ii) \quad \llbracket [\varphi, \psi] \rrbracket^{\mathcal{M}, g} = \llbracket \varphi \rrbracket^{\mathcal{M}, g} \times \llbracket \psi \rrbracket^{\mathcal{M}, g}$$

(ii*) where $\llbracket \varphi \rrbracket^{\mathcal{M}, g}$ and $\llbracket \psi \rrbracket^{\mathcal{M}, g}$ are defined

$$(iii) \quad \llbracket [\mathbf{X}: \varphi] \rrbracket^{\mathcal{M}, g} = \{f: \mathbf{X} \rightarrow \mathcal{U}_{\mathcal{M}} \mid \llbracket \varphi \rrbracket^{\mathcal{M}, g \cup f} = 1\}$$

(iii*) where $\text{dom}(g) \cap \mathbf{X} = \emptyset$ and $\llbracket \varphi \rrbracket^{\mathcal{M}, g \cup f}$ is defined

$$(iv) \quad \llbracket [\mathbf{K} \langle \mathbf{Q}x \rangle \mathbf{K}'] \rrbracket^{\mathcal{M}, g} = 1$$

iff $(\{u \in \mathcal{U}_{\mathcal{M}} \mid (\exists f) f \cup \{(x, u)\} \in \llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g}\},$

$$\{u \in \mathcal{U}_{\mathcal{M}} \mid (\exists f, h) [f \cup \{(x, u)\} \in \llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g} \ \& \ h \in \llbracket \mathbf{K}' \rrbracket^{\mathcal{M}, g \cup f \cup \{(x, u)\}}]\}) \in \mathcal{F}_{\mathcal{M}}(\mathbf{Q})$$

(iv*) where $\llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g}$ and $\llbracket \mathbf{K}' \rrbracket^{\mathcal{M}, g \cup f \cup \{(x, u)\}}$ are defined for all $f \in \llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g}$ and $u \in \mathcal{U}_{\mathcal{M}}$

Highly non-standard interpretation

Denotation $\llbracket \mathbf{A} \rrbracket^g$ of expression \mathbf{A} given model \mathcal{M} (as above) and $g: \text{Fr}(\mathbf{A}) \rightarrow \mathcal{U}_{\mathcal{M}}$

Notational convention: $\llbracket \mathbf{A} \rrbracket^{\mathcal{M}, \bar{g}} = \llbracket \mathbf{A} \rrbracket^{\mathcal{M}, g \upharpoonright \text{Fr}(\mathbf{A})}$

$$(i'') \quad \llbracket \mathbf{R}(x_1, \dots, x_n) \rrbracket^{\mathcal{M}, g} = \begin{cases} 1 & \text{iff } (g(x_1), \dots, g(x_n)) \in \mathcal{F}_{\mathcal{M}}(\mathbf{R}) \\ 0 & \text{iff } (g(x_1), \dots, g(x_n)) \notin \mathcal{F}_{\mathcal{M}}(\mathbf{R}) \end{cases} \quad [\approx (i')]$$

$$(ii'') \quad \llbracket [\varphi, \psi] \rrbracket^{\mathcal{M}, g} = \llbracket \varphi \rrbracket^{\mathcal{M}, \bar{g}} \times \llbracket \psi \rrbracket^{\mathcal{M}, \bar{g}} \quad [\approx (ii')]$$

$$(iii'') \quad \llbracket [\mathbf{X}: \varphi] \rrbracket^{\mathcal{M}, g} = \blacklozenge \{g: \mathbf{X} \rightarrow \mathcal{U}_{\mathcal{M}} \mid \llbracket \varphi \rrbracket^{\mathcal{M}, f \cup g} = 1\}$$

where $\blacklozenge R = \{h \text{ of } \mid h \in R \ \& \ f: \mid \text{dom}(h) \mid \xrightarrow{1-1} \text{onto} \text{dom}(h)\}$

$$(iv'') \quad \llbracket [\mathbf{K} \langle \mathbf{Q}x \rangle \mathbf{K}'] \rrbracket^{\mathcal{M}, g} = 1$$

iff $(\{u \in \mathcal{U}_{\mathcal{M}} \mid (\exists f) f[x/u] \in \bigcup \llbracket \mathbf{K} \rrbracket^{\mathcal{M}, \bar{g}}\},$

$[\approx (iv')]$

$$\{u \in \mathcal{U}_{\mathcal{M}} \mid (\exists h) [h[x/u] \in \bigcup \llbracket \mathbf{K} \oplus \mathbf{K}' \rrbracket^{\mathcal{M}, \bar{g}}]\}) \in \mathcal{F}_{\mathcal{M}}(\mathbf{Q})$$

non-compositional

where $[\mathbf{X}: \varphi] \oplus [\mathbf{Y}: \psi] = [\mathbf{X} \cup \mathbf{Y}: [\varphi, \psi]]$

Merge

and $g[x/u] = (g \setminus \{(x, g(x))\}) \cup \{(x, u)\}$

Conjecture

$$\llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g} = \blacklozenge \llbracket \mathbf{K} \rrbracket^{\mathcal{M}, g}, \text{ for any 'decent' DRS } \mathbf{K}, \text{ model } \mathcal{M}, \text{ and } g: \text{Fr}(\mathbf{A}) \rightarrow \mathcal{U}_{\mathcal{M}}$$

References

- Kamp, Hans: 'A Theory of Truth and Semantic Representation'. In: J. Groenendijk *et al.* (eds.), *Formal Methods in the Study of Language. Part 1*. Amsterdam 1981. 277–322.
 Kamp, Hans; Reyle, Uwe: *From Discourse to Logic*. Dordrecht 1993.